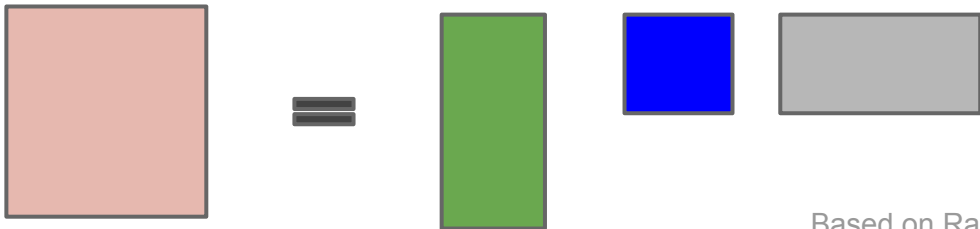
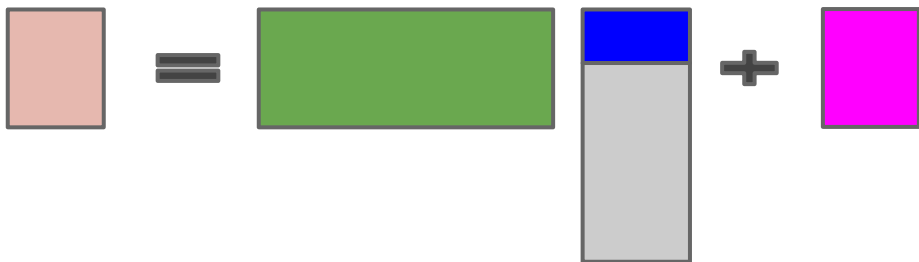


- 1. Greedy Algorithms for Structurally Constrained High Dimensional Problems**
- 2. Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization**

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CSci 8980

Motivation



Goal: Unifying
computational
framework for high-
dimensional
structured problems

Statistical Framework: Atomic Sets

Given a set of 'atoms' A

$$x = \sum_{i=1}^k c_i a_i, a_i \in A$$

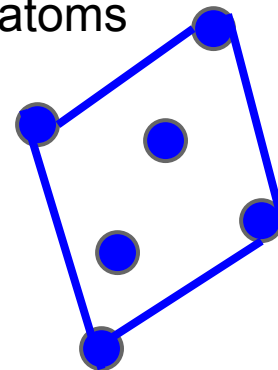
If C_A convex hull of A , **Gauge function:** $\|x\|_A = \inf\{t \geq 0 : x \in tC_A\}$

When A is centrosymmetric,

$$\|x\|_A = \inf \left\{ \sum_{a \in A} |c_a| : x = \sum_{a \in A} c_a \cdot a \right\}$$

Support function: $\|x\|_A^* = \sup \{ \langle x, a \rangle : a \in A \}$

D : convex hull of atoms



Any *linear* function will attain its *minimum* over D at an *atom* $s \in A$

Greedy algorithm (Tewari et al. 2011)

$$\min_{x: \|x\|_A \leq \kappa} f(x)$$

where f is **convex**, **smooth** and atomic norm is bounded $\{x : \|x\|_A \leq \kappa\}$ and f : goodness of fit measure.

Greedy Algorithm to minimize convex function f over κ scaled atomic norm ball

Let $x^0 = \kappa a_0$ for an arbitrary atom $a_0 \in A$
for $t = 0, \dots$, do

$$a_t = \arg \min_{a \in A} \langle \nabla f(x_t), a \rangle \quad (1)$$

(Line search for step size) $\gamma_t = \arg \min_{\gamma \in [0,1]} f(x_t + \gamma(\kappa a_t - x_t))$

(Update as linear combination) $x_{t+1} = x_t + \gamma_t(\kappa a_t - x_t)$

end for

- Add atom at every step
- Iterate x_t : conv. combination of at most $t+1$ atoms
- Select atom that makes the optimization problem easy

Contributions (Tewari et al. 2011)

Restricted Smoothness

in high dimensions not good smoothness constants

Restricted Smoothness Property

$$L_{\|\cdot\|}(f; S) = \sup_{x, y \in S, \alpha \in (0, 1)} \frac{f((1 - \alpha)x + \alpha y) - f(x) - \langle \nabla f(x), \alpha(y - x) \rangle}{\alpha^2 \|y - x\|^2}$$

$L_{\|\cdot\|}(f; S) \geq 0$ (since f is convex). How smooth in S with respect to $\|\cdot\|$
Connection with L-Lipschitz: $L_{\|\cdot\|}(f; S) \leq L$ (cares about smoothness just in S).

Connection with Hessian: $\|\nabla^2 f(z)\| \leq H$

Convergence results:

Convergence

$$f(x_T) - f(x^*) \leq \frac{8\kappa^2 L_{\|\cdot\|}(f; \kappa C_A) \|A\|^2}{T}$$

$$\|A\| = \sup_{a \in A} \|a\|$$

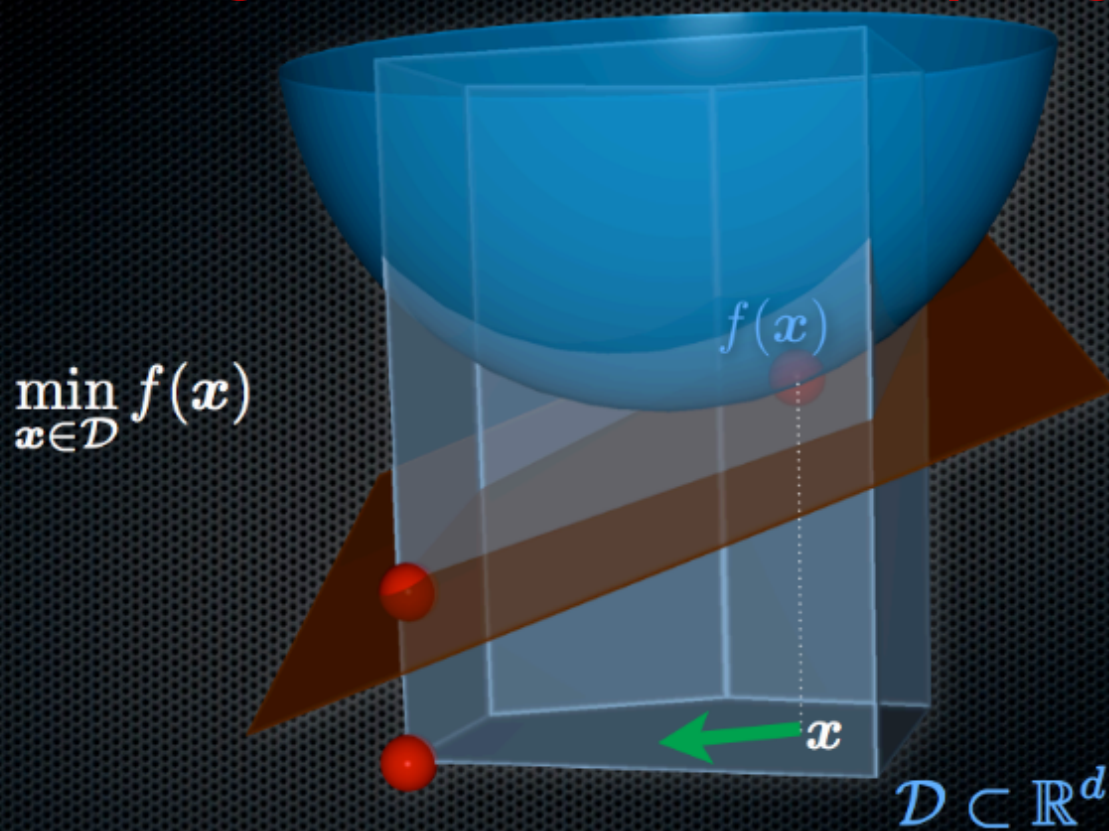
For Banach spaces:

Extension to infinite dimensional Banach spaces

V : Banach space equipped with inner product. ∇f : Fechel derivative: elements of the dual space V^* , inner product $\langle X, x \rangle = X(x)$, $x \in V, X \in V^*$

$$L_{\|\cdot\|}(f; S) = \sup_{x, y \in S, \alpha \in (0, 1)} \frac{f((1 - \alpha)x + \alpha y) - f(x) - \langle \nabla f(x), \alpha(y - x) \rangle}{\frac{1}{r} \alpha^r \|y - x\|^r}, r \in [1, 2]$$
$$\rightarrow O\left(\frac{1}{T^{r-1}}\right)$$

Revisiting Frank - Wolfe (Jaggi 2013)



Frank - Wolfe

Algorithm 1 Frank-Wolfe (1956)

Let $x^0 \in D$

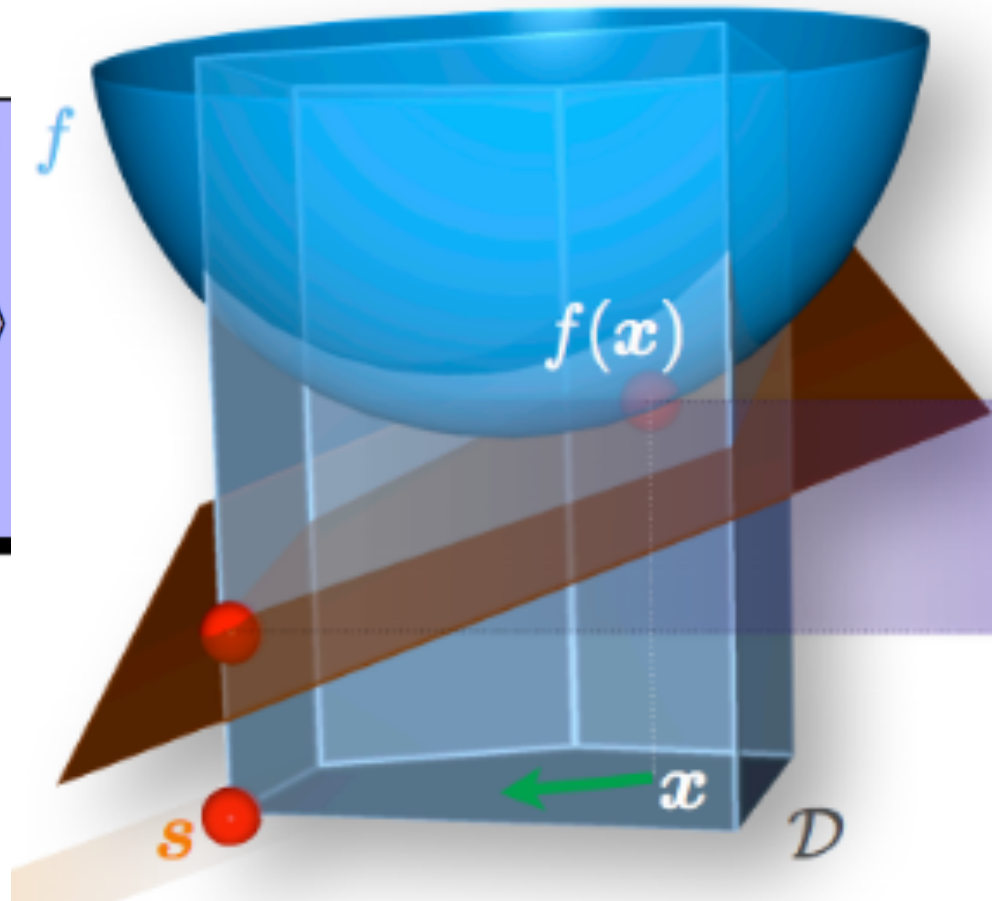
for $k = 0, \dots, K$ **do**

Compute $s = \arg \min_{s' \in D} \langle s', \nabla f(x^{(k)}) \rangle$

Let $\gamma = \frac{2}{k+2}$

Update $x^{(k+1)} = (1 - \gamma)x^{(k)} + \gamma s$

end for



The Duality gap and Certificates

Primal problem:

$$\min_{x \in D} f(x)$$

Surrogate duality gap.

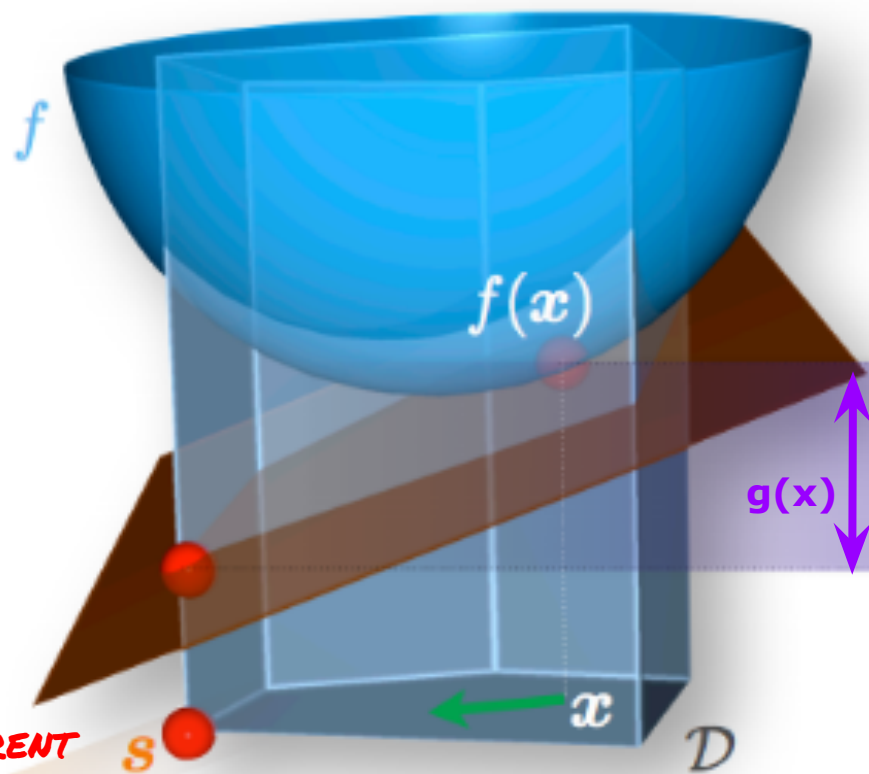
$$g(x) = \max_{s' \in D} \langle x - s', \nabla f(x) \rangle$$

Linearized problem

$$w(x) = \min_{s' \in D} f(x) + \langle s' - x, \nabla f(x) \rangle$$

$$g(x) \geq f(x) - f(x^*)$$

**CERTIFICATE FOR CURRENT
APPROXIMATION QUALITY**



Frank-Wolfe Variants

Algorithm 1 Frank-Wolfe (1956)

Let $x^0 \in D$
for $k = 0, \dots, K$ do
 Compute $s = \arg \min_{s' \in D} \langle s', \nabla f(x^{(k)}) \rangle$
 Let $\gamma = \frac{2}{k+2}$
 Update $x^{(k+1)} = (1 - \gamma)x^{(k)} + \gamma s$
end for

Algorithm 2 Frank-Wolfe with Approximate Linear Subproblems, for Quality $\delta \geq 0$

Let $x^0 \in D$
for $k = 0, \dots, K$ do
 Let $\gamma = \frac{2}{k+2}$
 Find $s \in D$ s.t

$$\langle s, \nabla f(x^{(k)}) \rangle \leq \min_{s' \in D} \langle s', \nabla f(x^{(k)}) \rangle + \frac{1}{2} \delta \gamma C_f$$

a) Optionally: perform line search for γ
b) Update $x^{(k+1)} = (1 - \gamma)x^{(k)} + \gamma s$
end for

Algorithm 3 Line-Search for the step size γ

As Algorithm 2, except replacing line a) with :

$$\gamma = \arg \min_{\gamma \in [0,1]} f(x^{(k)} + \gamma(s - x^{(k)}))$$

Algorithm 4 Fully Corrective, Re-optimizing over all previous directions

As Algorithm 2, except replacing line b) with :
Do the update

$$x^{(k+1)} = \arg \min_{x \in \text{conv}(s^{(0)}, \dots, s^{(k+1)})} f(x)$$

Contributions (Jaggi 2013)

The Curvature

$$C_f := \sup_{x, s \in D, \gamma \in [0, 1], y = x + \gamma(s-x)} \frac{2}{\gamma^2} (f(y) - f(x) - \langle y - x, \nabla f(x) \rangle)$$

Connection with L-Lipschitz: $C_f \leq \text{diam}_{\|\cdot\|}(D^2)L$

Convergence results:

Primal Convergence

$$f(x^{(k)}) - f(x^*) \leq \frac{2C_f}{k+2}(1+\delta)$$

δ accuracy for solving the linear sub-problems.

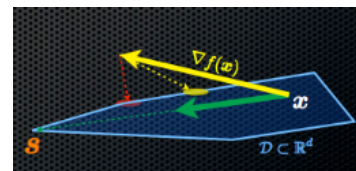
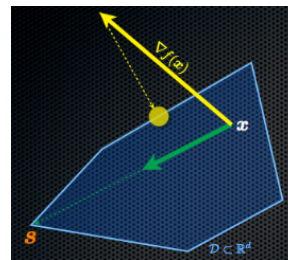
Primal-Dual Convergence

$$g(x^{(\hat{k})}) \leq \frac{7C_f}{K+2}(1+\delta)$$

where $1 \leq \hat{k} \leq K$

CONTRIBUTION NO.1

1. Duality gap convergence guarantee
2. Affine invariance



3. Optimality in terms of sparsity

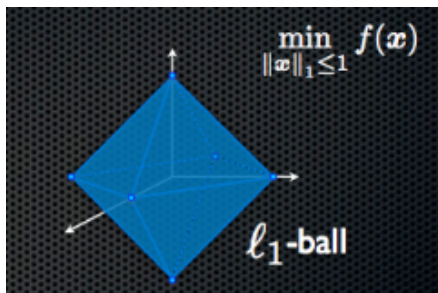
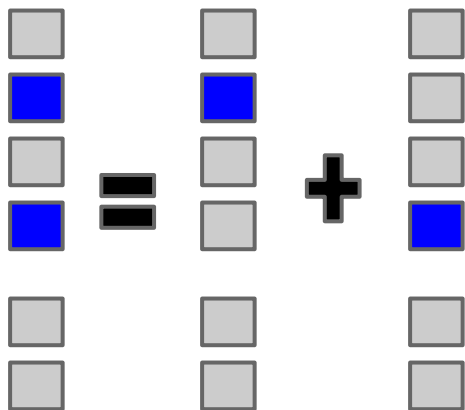
The obtained sparsity k is **optimal** for an approximation quality of $1/k$

Frank - Wolfe on Atomic Domains

\mathcal{X}	Optimization Domain		Complexity of one Frank-Wolfe Iteration	
	Atoms \mathcal{A}	$\mathcal{D} = \text{conv}(\mathcal{A})$	$\sup_{s \in \mathcal{D}} \langle s, \mathbf{y} \rangle$	Complexity
\mathbb{R}^n	Sparse Vectors	$\ \cdot\ _1$ -ball	$\ \mathbf{y}\ _\infty$	$O(n)$
\mathbb{R}^n	Sign-Vectors	$\ \cdot\ _\infty$ -ball	$\ \mathbf{y}\ _1$	$O(n)$
\mathbb{R}^n	ℓ_p -Sphere	$\ \cdot\ _p$ -ball	$\ \mathbf{y}\ _q$	$O(n)$
\mathbb{R}^n	Sparse Non-neg. Vectors	Simplex Δ_n	$\max_i \{y_i\}$	$O(n)$
\mathbb{R}^n	Latent Group Sparse Vec.	$\ \cdot\ _{\mathcal{G}}$ -ball	$\max_{g \in \mathcal{G}} \ \mathbf{y}_{(g)}\ _q^*$	$\sum_{g \in \mathcal{G}} g $
$\mathbb{R}^{m \times n}$	Matrix Trace Norm	$\ \cdot\ _{tr}$ -ball	$\ \mathbf{y}\ _{op} = \sigma_1(\mathbf{y})$	$\tilde{O}(N_f/\sqrt{\epsilon'})$ (Lanczos)
$\mathbb{R}^{m \times n}$	Matrix Operator Norm	$\ \cdot\ _{op}$ -ball	$\ \mathbf{y}\ _{tr} = \ (\sigma_i(\mathbf{y}))\ _1$	SVD
$\mathbb{R}^{m \times n}$	Schatten Matrix Norms	$\ (\sigma_i(\cdot))\ _p$ -ball	$\ (\sigma_i(\mathbf{y}))\ _q$	SVD
$\mathbb{R}^{m \times n}$	Matrix Max-Norm	$\ \cdot\ _{\max}$ -ball		$\tilde{O}(N_f(n+m)^{1.5}/\epsilon'^{2.5})$
$\mathbb{R}^{n \times n}$	Permutation Matrices	Birkhoff polytope		$O(n^3)$
$\mathbb{R}^{n \times n}$	Rotation Matrices			SVD (Procrustes prob.)
$\mathbb{S}^{n \times n}$	Rank-1 PSD matrices of unit trace	$\{x \succeq 0, \text{Tr}(x)=1\}$	$\lambda_{\max}(\mathbf{y})$	$\tilde{O}(N_f/\sqrt{\epsilon'})$ (Lanczos)
$\mathbb{S}^{n \times n}$	PSD matrices of bounded diagonal	$\{x \succeq 0, x_{ii} \leq 1\}$		$\tilde{O}(N_f n^{1.5}/\epsilon'^{2.5})$

Table: from Jaggi's presentation in Smile, Paris Seminar, 2013

Special case: Sparse vectors



Set of atoms: $A = \{\pm e_i | i \in [n]\}$

$\text{Conv}(A) =$ unit ball of ℓ_1 norm

Compute

$$a_t = \arg \min_{a \in \pm\{e_1, \dots, e_p\}} \langle \nabla f(x_t), a \rangle$$

Greedy coordinate descent: Find $j = \arg \max_{j' \in \{1, \dots, p\}} |[\nabla f(x_t)]_{j'}|$
and set $a_t = -\text{sign}([\nabla f(x_t)]_j) e_j$

Obtain $O(1/k)$ approximate solution of sparsity k
Equivalent to (Orthogonal) Matching Pursuit.
Original Frank-Wolfe to polyhedral sets.

Special case: sparse non-negative vectors

Set of atoms: $A = \{e_i | i \in [n]\}$

$\text{Conv}(A) = \text{Simplex}$

Compute

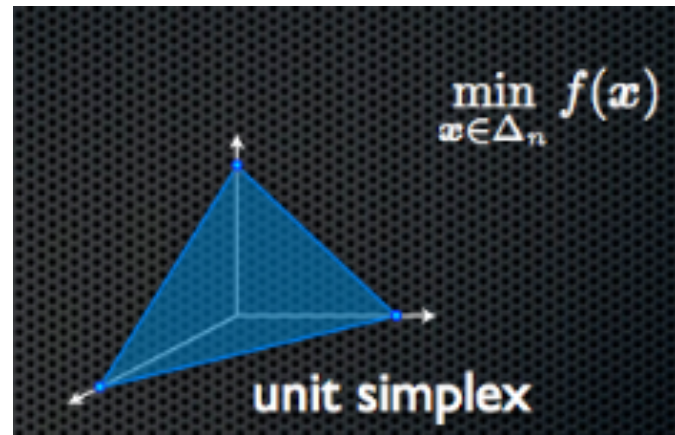
$$a_t = \arg \min_{a \in \pm\{e_1, \dots, e_p\}} \langle \nabla f(x_t), a \rangle$$

Find $j = \arg \min_{j' \in \{1, \dots, p\}} [\nabla f(x_t)]_{j'}$ and set $a_t = e_j$

Clarkson 2010.

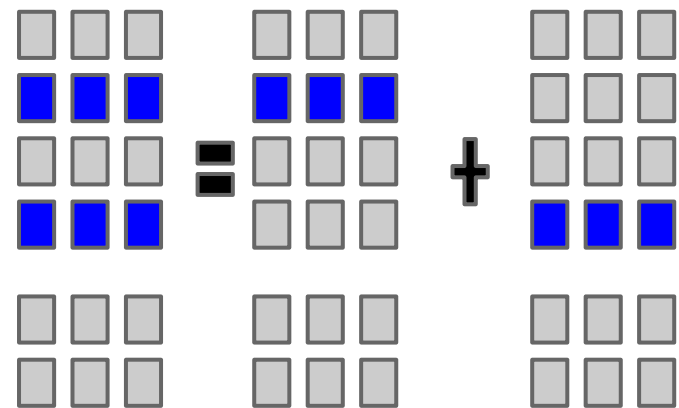
Tradeoff between **sparsity** and **approximation quality**

Sparsity of FW iterates is **Optimal** in primal and dual approximation quality



If we choose $\| \cdot \|$ to be l_1 norm, then restricted smoothness constant is similar to C_f

Special case: Group Sparse Matrices



Infinite Set of atoms: All matrices with a single non-zero row where the row has ℓ_q norm 1 for some $q > 1$

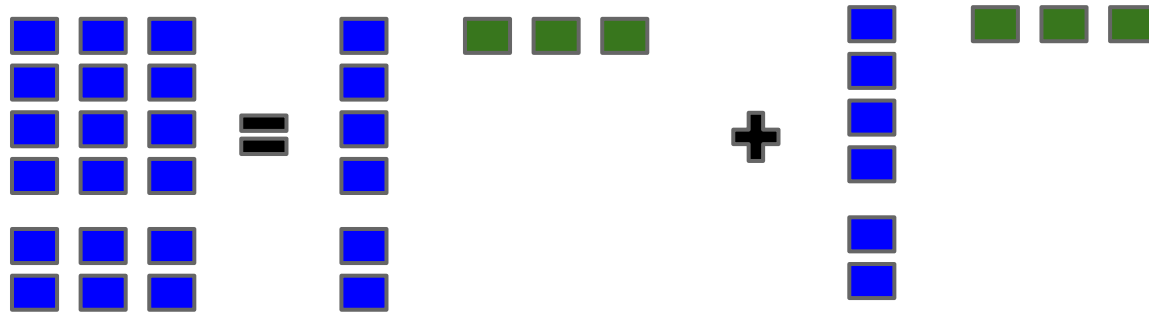
$\text{Conv}(A) =$ unit ball of the $\|\cdot\|_{q,1}$ group norm on $R^{p \times k}$ (sum of ℓ_q norms on the rows)

Compute

$$a_t = \arg \min_{a: \text{nnzrows}(a)=1, \|a\|_{q,1}=1} \langle \nabla f(x_t), a \rangle$$

Find row j of $\nabla f(x_t)$ with maximal $\ell_{q'}$ norm. Set a_t to be the matrix with the only non zero row in j equal to u^T such that $\langle u, [\nabla f(x_t)]_{j,:}^T \rangle = -\|[\nabla f(x_t)]_{j,:}^T\|_{q'}$

Special case: Low rank matrices



Infinite Set of atoms: All rank 1 matrices with Frobenius norm = 1

$$A := \{uv^T \mid u \in \mathbb{R}^n, \|u\|_2 = 1, v \in \mathbb{R}^n, \|v\|_2 = 1\}$$

Conv(A) = unit ball of the Trace norm (Schatten ℓ_1 norm)

Greedy step: Compute

$$a_t = \arg \min_{a: \text{rank}(a)=1, \|a\|_F=1} \langle \nabla f(x_t), a \rangle$$

Compute SVD $\nabla f(x_t) = U\Sigma V^T$ and set $a = -u_1 v_1^T$, where u_1, v_1 left and right singular vectors corresponding to largest singular value.

Polynomial
time

- ❑ (Tewari et al 2011) Not polynomial time to compute greedy step for **non negative low rank matrices**
- ❑ **Permutation matrices:** Efficient optimization over Birkoff polytope, Hungarian Algorithm (Conv(A) = set of doubly stochastic matrices)
- ❑ (Jaggi 2013)

Novel framework for factorized matrix norms

Novel framework for matrix factorization:

$$A = \{LR^T \mid L \in A_{\text{left}}, R \in A_{\text{right}}\},$$

$A_{\text{left}} \subseteq R^{m \times r}, A_{\text{right}} \subseteq R^{n \times r}$
All Frank-Wolfe iterates are low-rank updates: $s = LR^T$

Thanks

References

- 1** Greedy Algorithms for Structurally Constrained High Dimensional Problems, A. Tewari, P. Ravikumar, I. Dhillon. In Advances in Neural Information Processing Systems (NIPS) 24, 2011.
- 2** Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization, ICML 2013